

Assessing The (In)Ability of LLMs To Reason in Interval Temporal Logic

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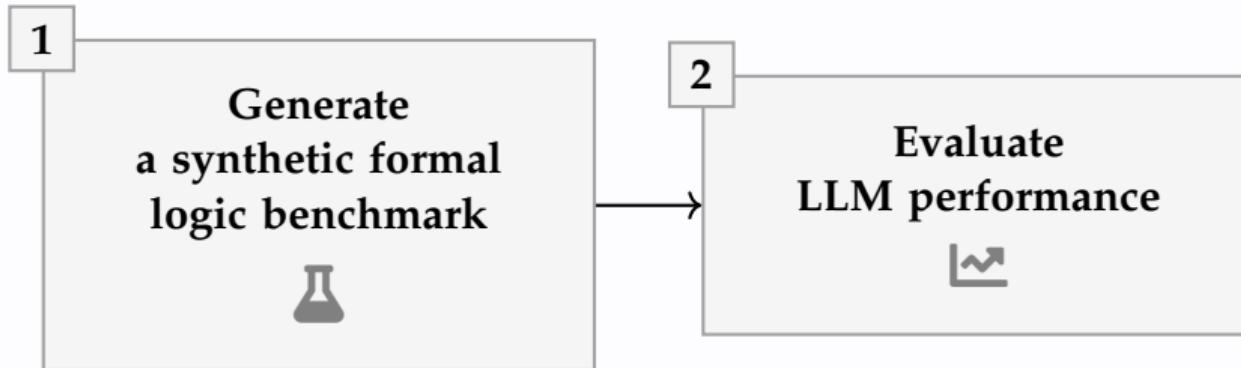


I

Introduction



Objective



Why synthetic benchmarks?



Contamination

Often solutions become part of the training corpus of LLMs.

Training data decontamination procedures are often only partially effective.



Scalability

Manually created benchmarks are difficult to scale both in *size* and *complexity*, requiring significant human effort and financial investments.

II.I

Generation (Logic)



Halpern–Shoham’s Logic

Halpern–Shoham’s Logic (HS) is a modal logic where *intervals*, rather than points, are the fundamental states.

Key idea: Accessibility between intervals is determined by *Allen’s interval relations*.

Given a linear order $\mathbb{D} = \langle D, < \rangle$, a *strict interval* is an ordered pair $[x, y]$ with $x, y \in D$ and $x < y$.

Two intervals $[x, y]$ and $[w, z]$ are compared by their endpoints, and their relation is captured by one of Allen’s modalities.

HS introduces an *existential modality* $\langle X \rangle$ for each Allen relation R_X . The six basic relations A, L, B, E, D, O each have an inverse \bar{X} , yielding 12 binary relations in total.

Interval Relations and Modalities

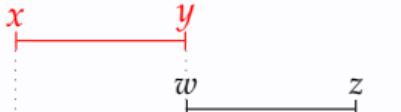
HS modality	Definition w.r.t. the interval structure	Example
$\langle A \rangle$ (adjacent)	$[x, y]R_A[w, z] \Leftrightarrow y = w$	
$\langle L \rangle$ (later)	$[x, y]R_L[w, z] \Leftrightarrow y < w$	
$\langle B \rangle$ (begins)	$[x, y]R_B[w, z] \Leftrightarrow x = w \wedge z < y$	
$\langle E \rangle$ (ends)	$[x, y]R_E[w, z] \Leftrightarrow y = z \wedge x < w$	
$\langle D \rangle$ (during)	$[x, y]R_D[w, z] \Leftrightarrow x < w \wedge z < y$	
$\langle O \rangle$ (overlaps)	$[x, y]R_O[w, z] \Leftrightarrow x < w < y < z$	

Table: Allen's interval relations and HS modalities.

Syntax of HS

Alphabet: propositional letters \mathcal{P} , classical connectives \neg, \vee , and modalities $\langle X \rangle$ for $X \in \mathcal{X}$ with:

$$\mathcal{X} = \{A, \overline{A}, L, \overline{L}, B, \overline{B}, E, \overline{E}, D, \overline{D}, O, \overline{O}\}.$$

Grammar:

$$\varphi ::= p \mid \neg\varphi \mid \varphi \vee \varphi \mid \langle X \rangle \varphi \quad (p \in \mathcal{P}, X \in \mathcal{X}).$$

Derived connectives: $\varphi \wedge \psi \equiv \neg(\neg\varphi \vee \neg\psi)$, $\varphi \rightarrow \psi \equiv \neg\varphi \vee \psi$, $\top \equiv p \vee \neg p$.

Universal modality: $[X]\varphi \equiv \neg\langle X \rangle \neg\varphi$.

Interval model: $M = \langle \mathbb{I}(\mathbb{D}), V \rangle$ with \mathbb{D} a linear order, $\mathbb{I}(\mathbb{D})$ the set of strict intervals over \mathbb{D} , and $V : \mathcal{P} \rightarrow 2^{\mathbb{I}(\mathbb{D})}$ a valuation.

Truth on an interval $[x, y]$:

$$M, [x, y] \models p \Leftrightarrow [x, y] \in V(p) \quad (p \in \mathcal{P}),$$

$$M, [x, y] \models \neg\psi \Leftrightarrow M, [x, y] \not\models \psi,$$

$$M, [x, y] \models \psi_1 \vee \psi_2 \Leftrightarrow M, [x, y] \models \psi_1 \text{ or } M, [x, y] \models \psi_2,$$

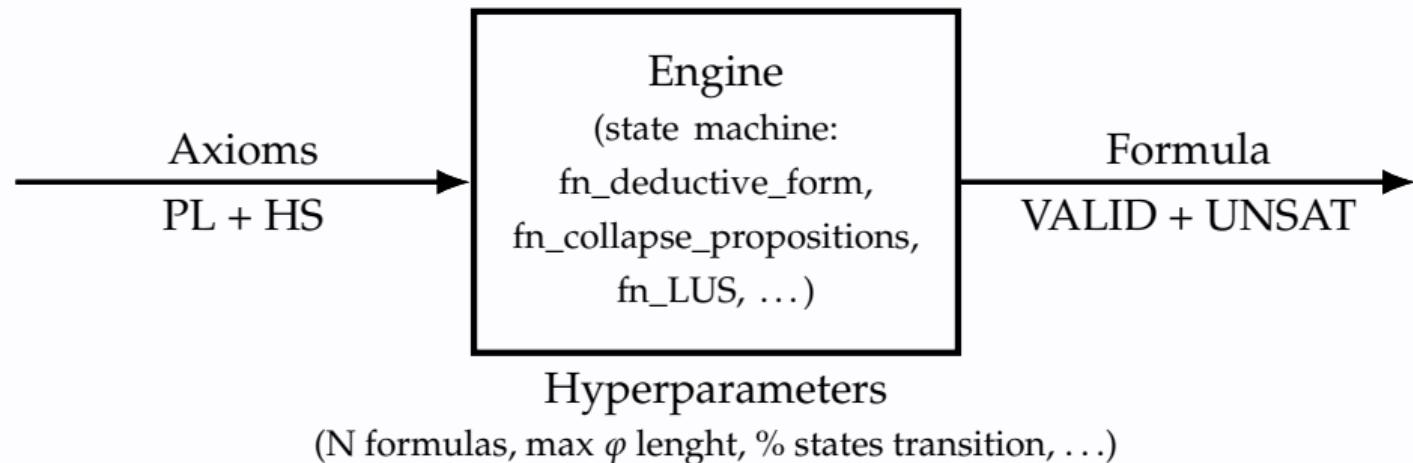
$$M, [x, y] \models \langle X \rangle \psi \Leftrightarrow \exists [w, z] \text{ s.t. } [x, y] R_X [w, z] \text{ and } M, [w, z] \models \psi.$$

II.II

Generation (Scheme)

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Scheme



II.III

Generation (Input)

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PL Axioms

Axiom	Type
$(a \rightarrow b) \leftrightarrow (\neg b \rightarrow \neg a)$	contrapositive
$(a \rightarrow b) \leftrightarrow (\neg a \vee b)$	implication as disjunction
$a \wedge (b \vee c) \leftrightarrow (a \wedge b) \vee (a \wedge c)$	distributivity of \wedge over \vee
$\neg(a \wedge b) \leftrightarrow \neg a \vee \neg b$	De Morgan
$\neg(a \vee b) \leftrightarrow \neg a \wedge \neg b$	De Morgan
$a \wedge a \leftrightarrow a$	idempotence
$a \wedge \top \leftrightarrow a$	identity
$\neg \neg a \leftrightarrow a$	double negation
$a \wedge (a \rightarrow b) \rightarrow b$	modus ponens
$(a \rightarrow b) \wedge \neg b \rightarrow \neg a$	modus tollens
$\neg a \wedge (a \vee b) \rightarrow b$	disjunctive syllogism

Table: All propositional axioms used.

HS Axioms

Axiom	Type
$\langle \bar{B} \rangle \langle \bar{E} \rangle p \leftrightarrow \langle \bar{E} \rangle \langle \bar{B} \rangle p$	commutativity
$\langle L \rangle p \leftrightarrow \langle A \rangle \langle A \rangle p$	definability
$\neg [B]p \leftrightarrow \langle B \rangle \neg p$	duality
$[A](p \rightarrow q) \rightarrow ([A]p \rightarrow [A]q)$	K axiom
$\langle B \rangle \langle B \rangle p \rightarrow \langle B \rangle p$	transitivity
$\langle B \rangle [\bar{B}]p \rightarrow p$	temporality
$p \rightarrow [A] \langle \bar{A} \rangle p$	inverse of temporality
$\langle A \rangle \langle \bar{A} \rangle p \rightarrow [A] \langle \bar{A} \rangle p$	stability

Table: Selected HS axioms, one for each type.

II.IV

Generation (Output)

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Output

id_enriched	premise_size	conclusion_size	total_size	compositional_hops	id_base	id_plus
Te1	6	7	14	1	Ae9	Ae41
					base: $\neg(t \vee q) \models (\neg t \wedge \neg q)$ plus: $\neg[\text{later}](s) \models \langle \text{later} \rangle \neg s$ enriched: $\neg(t \vee \neg[\text{later}](s)) \models (\neg t \wedge \neg\langle \text{later} \rangle \neg s)$ enriched_valid: $\neg(\neg(t \vee \neg[\text{later}](s))) \wedge (t \vee \langle \text{later} \rangle \neg s)$ enriched_unsat: $\neg(t \vee \neg[\text{later}](s)) \wedge (\neg(\neg t) \vee \langle \text{later} \rangle \neg s)$	
Ti8	7	7	15	2	Ti3	Ae54
					base: $\langle \text{meets} \rangle \langle \text{met_by} \rangle \neg[\text{begins}](q) \models [\text{meets}] \langle \text{met_by} \rangle \langle \text{begins} \rangle \neg q$ plus: $\langle \text{during} \rangle \neg p \models \neg[\text{during}](p)$ enriched: $\langle \text{meets} \rangle \langle \text{met_by} \rangle \neg[\text{begins}] (\langle \text{during} \rangle \neg p) \models [\text{meets}] \langle \text{met_by} \rangle \langle \text{begins} \rangle (\neg \neg[\text{during}](p))$ enriched_valid: $[\text{meets}] [\text{met_by}] [\text{begins}] (\langle \text{during} \rangle \neg p) \vee [\text{meets}] \langle \text{met_by} \rangle \langle \text{begins} \rangle [\text{during}](p)$ enriched_unsat: $(\langle \text{meets} \rangle \langle \text{met_by} \rangle \neg[\text{begins}] (\langle \text{during} \rangle \neg p)) \wedge \langle \text{meets} \rangle [\text{met_by}] \langle \text{begins} \rangle \neg[\text{during}](p)$	

Table: Example of raw output of two sets of valid and unsatisfiable formulas.

II.V

Generation (Engine)

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Formulas generation

Formula Structures

Equivalence form: $\varphi \leftrightarrow \psi$

Implicative form: $\varphi \rightarrow \psi$

Formula Roles

Base types: equivalence or implicative

Plus types: only equivalence

Substitution Principle

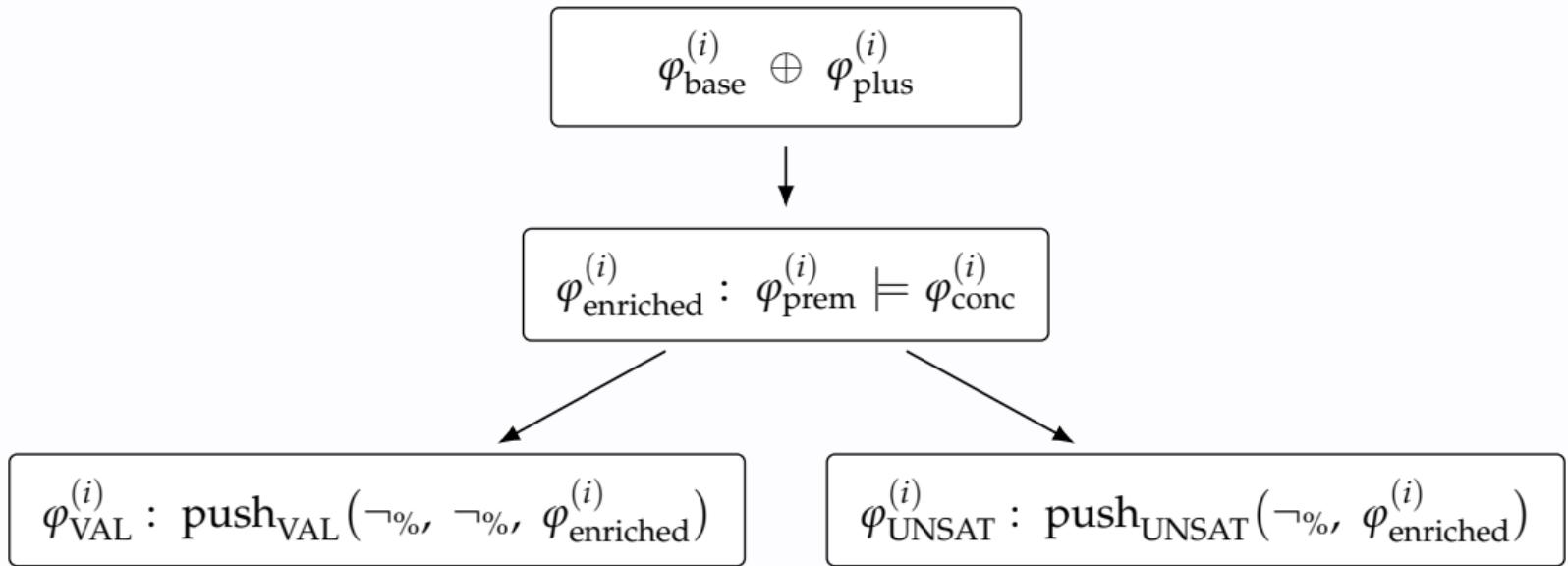
Inspired by uniform substitution:

$$\varphi_{prem}^{base}[p/\psi] \models \varphi_{conc}^{base}[p/\psi]$$

Instead of a WFF ψ , we use a plus formula:

$$\varphi_{enriched} : \varphi_{prem}^{base}[p/\varphi_{prem}^{plus}]_{partial} \models \varphi_{conc}^{base}[p/\varphi_{conc}^{plus}]$$

Valid and unsatisfiable formulas



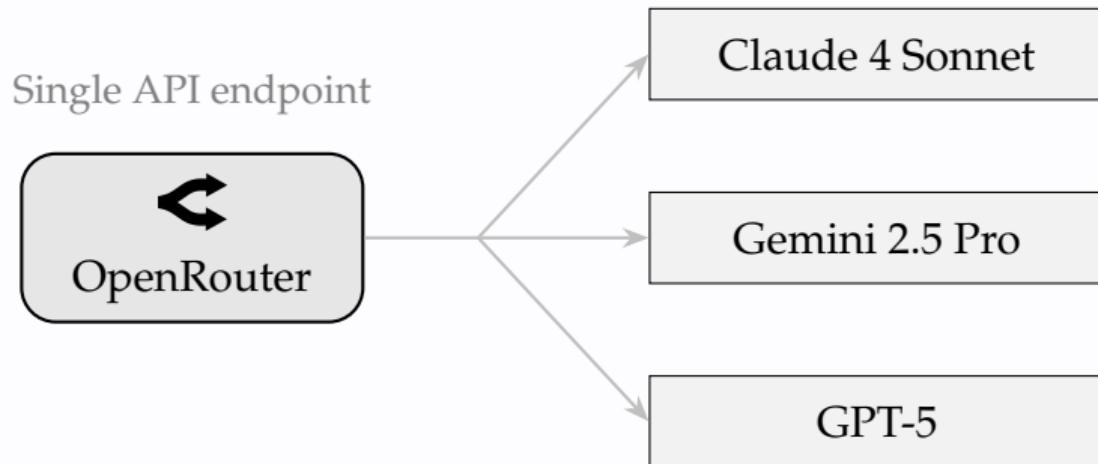
Tree view of Formula Construction

```
Ti77: ([finished](!(<begun_by>(!((p & !([before](t))))))) | //
|     // [met_by](<meets>(<overlaps>((!(p) | !(<before>(!!(t)))))))
|-- Ti32b: ([finished]([begun_by]((p & t))) | //
|     // [met_by](<meets>(<overlaps>((!(p) | !(t))))))
|     |-- Ti18b: ([finished]([begun_by](!!(p))) | //
|     |     // [met_by](<meets>(<overlaps>(p)))
|     |     |-- Ti15b: (!!(q) | [met_by](<meets>(q)))
|     |     |     |-- Ai38b: p |= [met_by](<meets>(p))
|     |     |     |-- Ae16p: q |= !(!!(q))
|     |     |     |-- Ae30p: <finished>(<begun_by>(p)) |= <overlaps>(p)
|     |     |-- Ae7p: !((p & t)) |= (!!(p) | !(t))
|-- Ae43p: !([before](t)) |= <before>(!!(t))
```

III

Evaluation





Adopted prompting strategies



Few-Shot

Provides **complete examples** of problems with their solutions to facilitate learning by analogy without explicit instructions.



CoT

Guides the model to **decompose** the problem into multiple components to facilitate its resolution.



Context

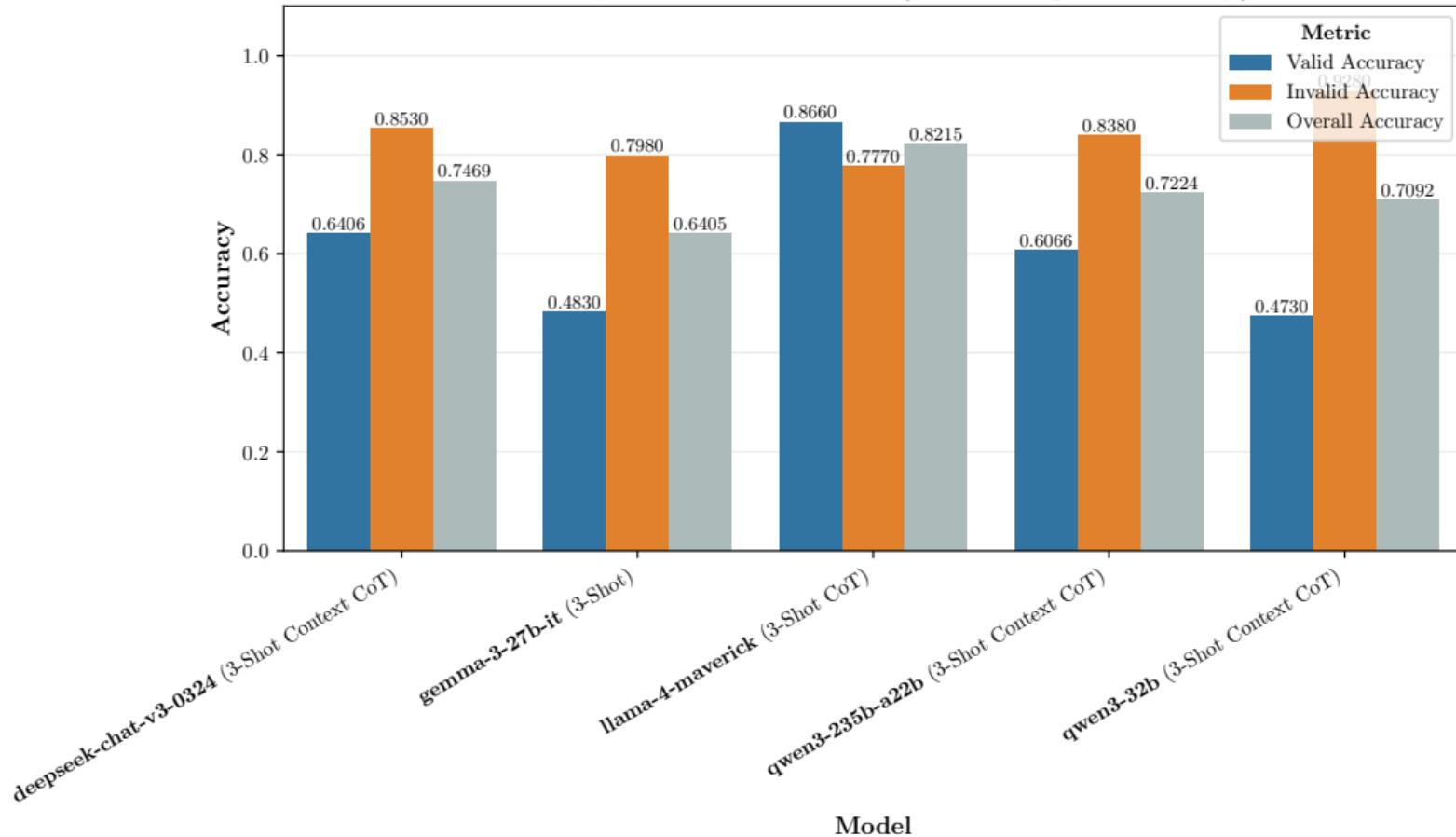
A natural language **introduction** to LTL is provided followed by its syntax and semantics. A natural language **introduction** to HS is provided followed by its syntax and semantics.

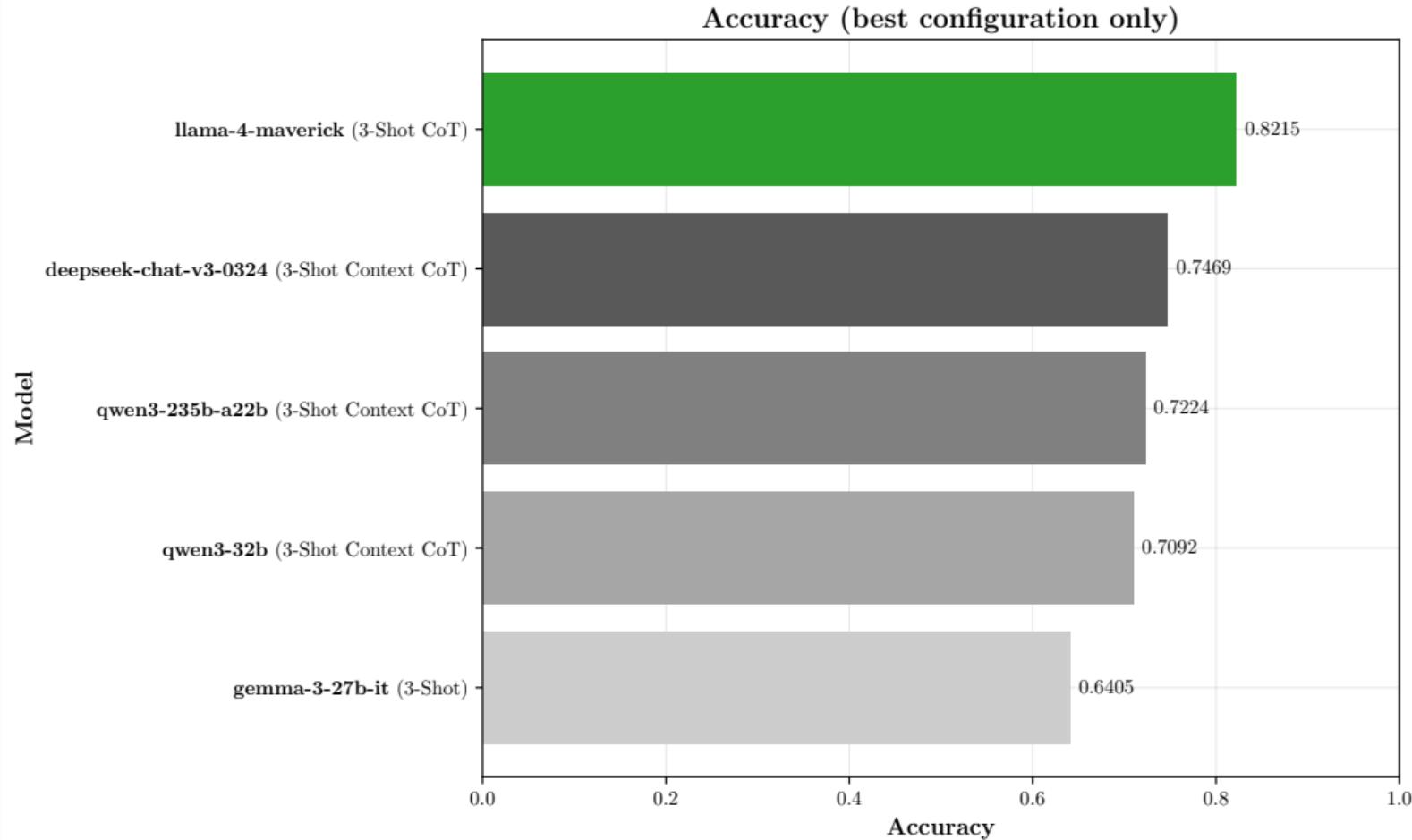
IV

Results

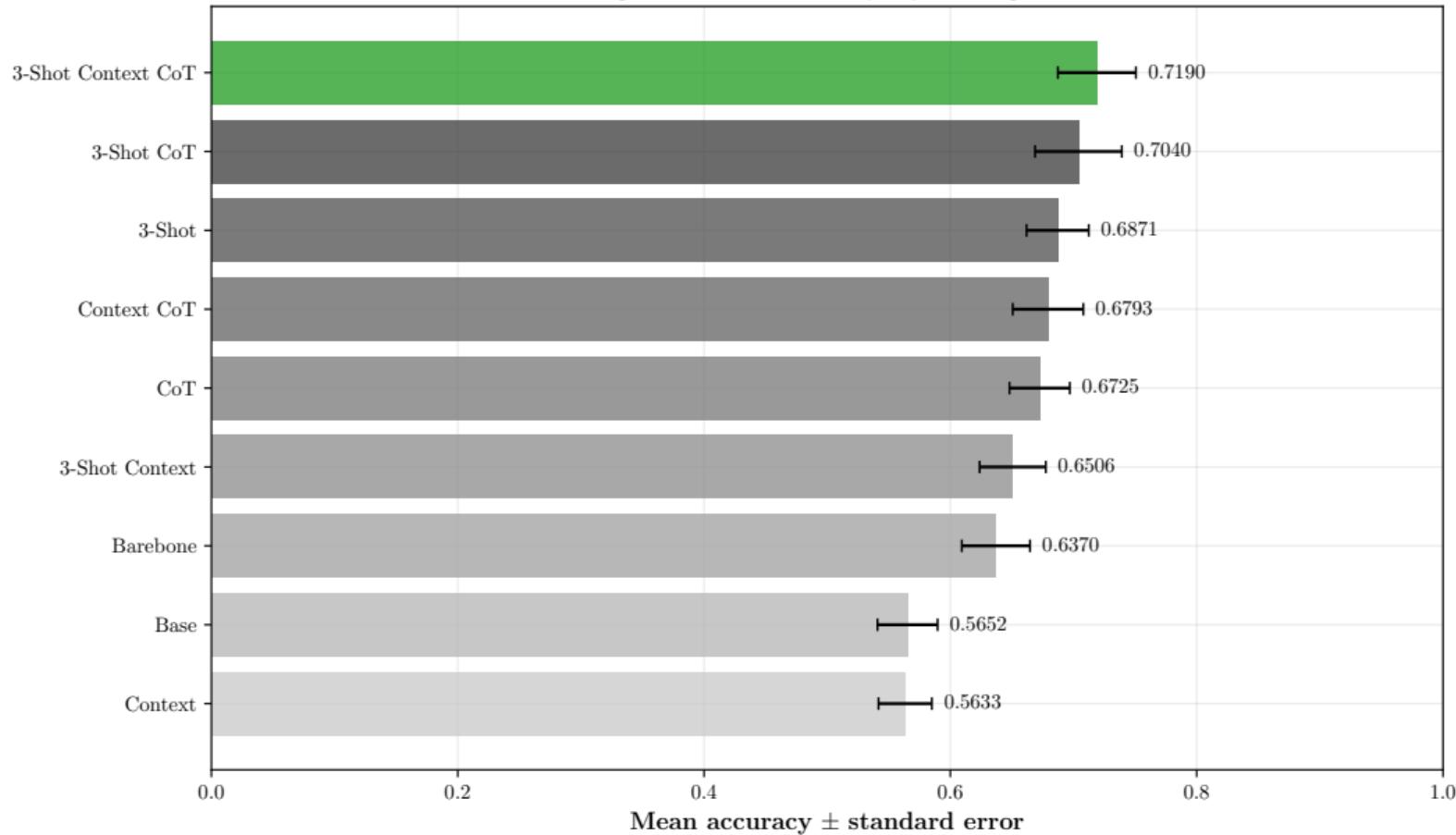


Model performance metrics (best configuration only)





Average overall accuracy by configuration



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Appendix

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Overall accuracy for all model configurations

